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Low-Depth Spatial Tree Algorithms

IPDPS 2024



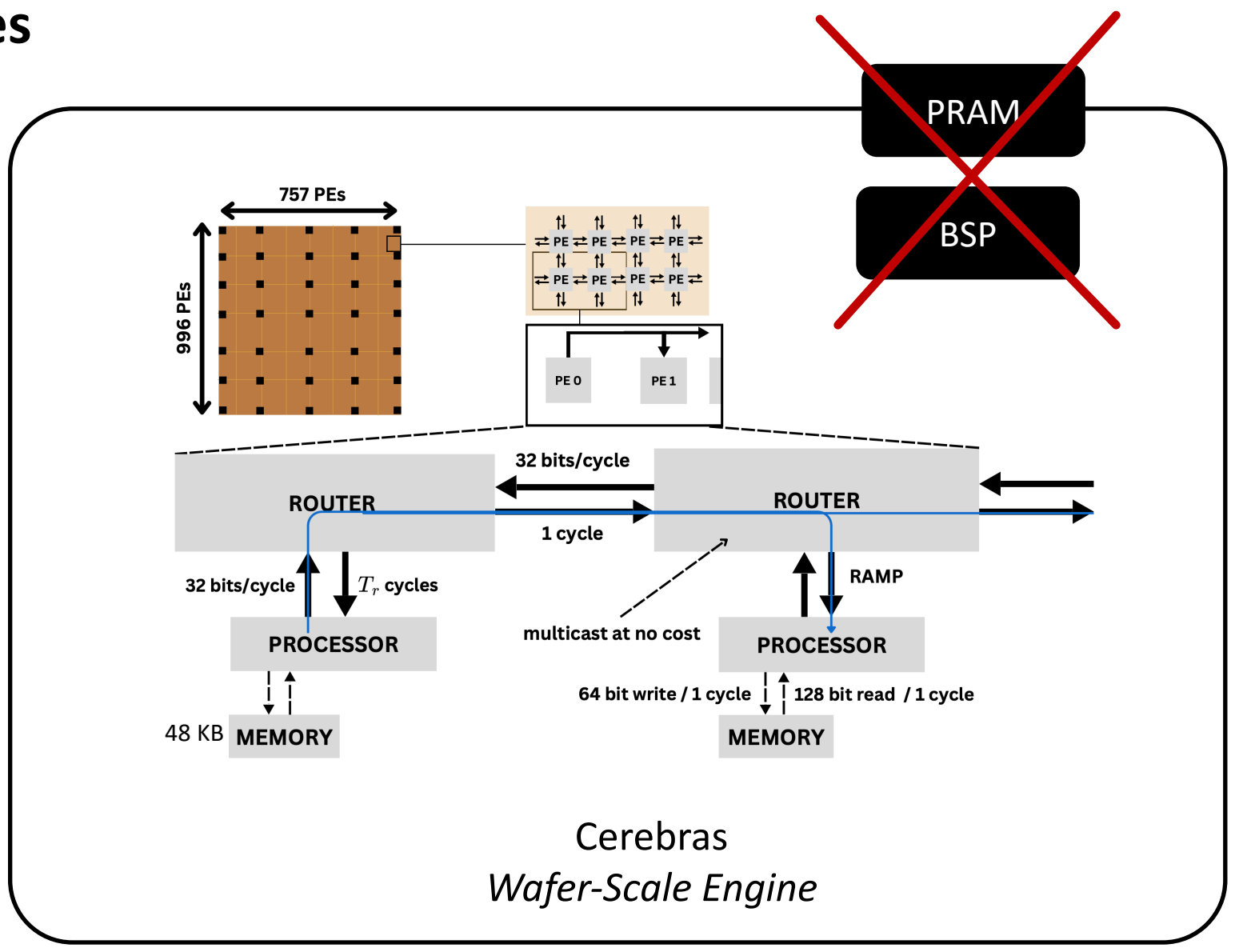
Spatial Dataflow Architectures

AMD
Versal Adaptive SoCs

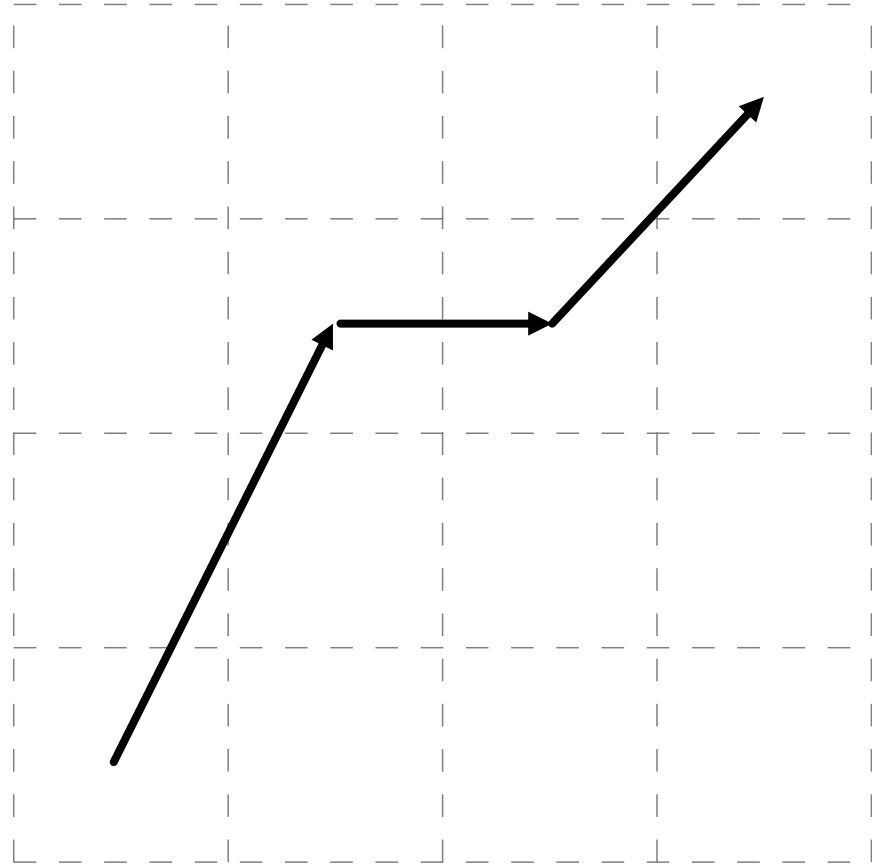
Samba Nova
Reconfigurable Dataflow Architecture

...

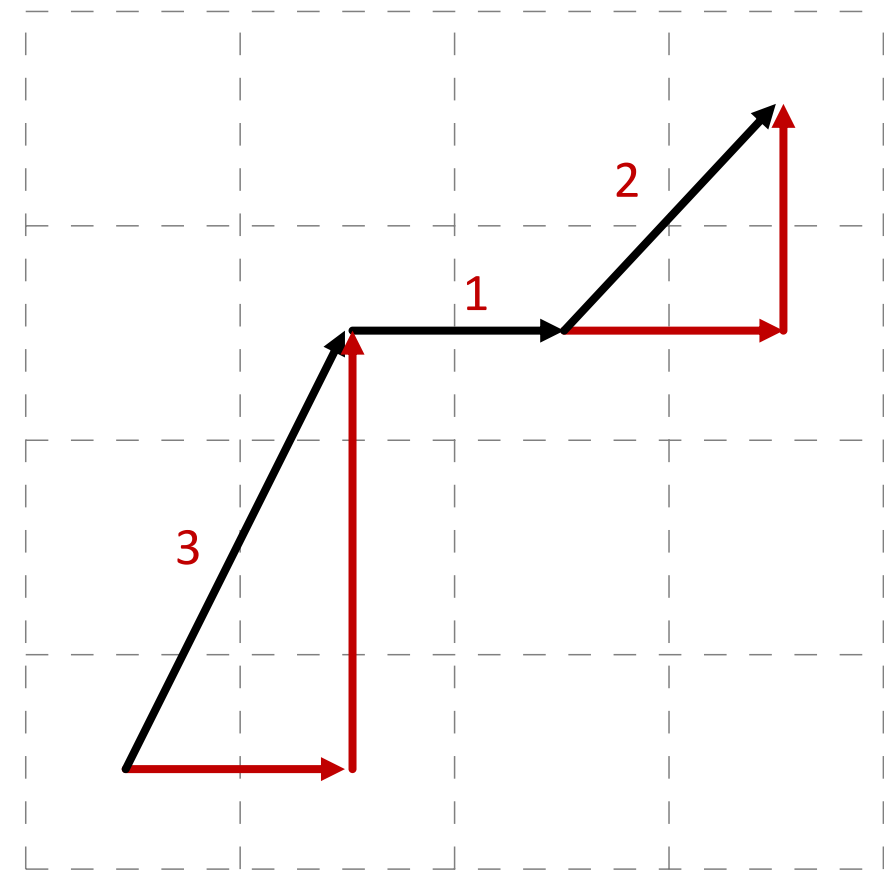
Graphcore
IPUs



Model of Computation

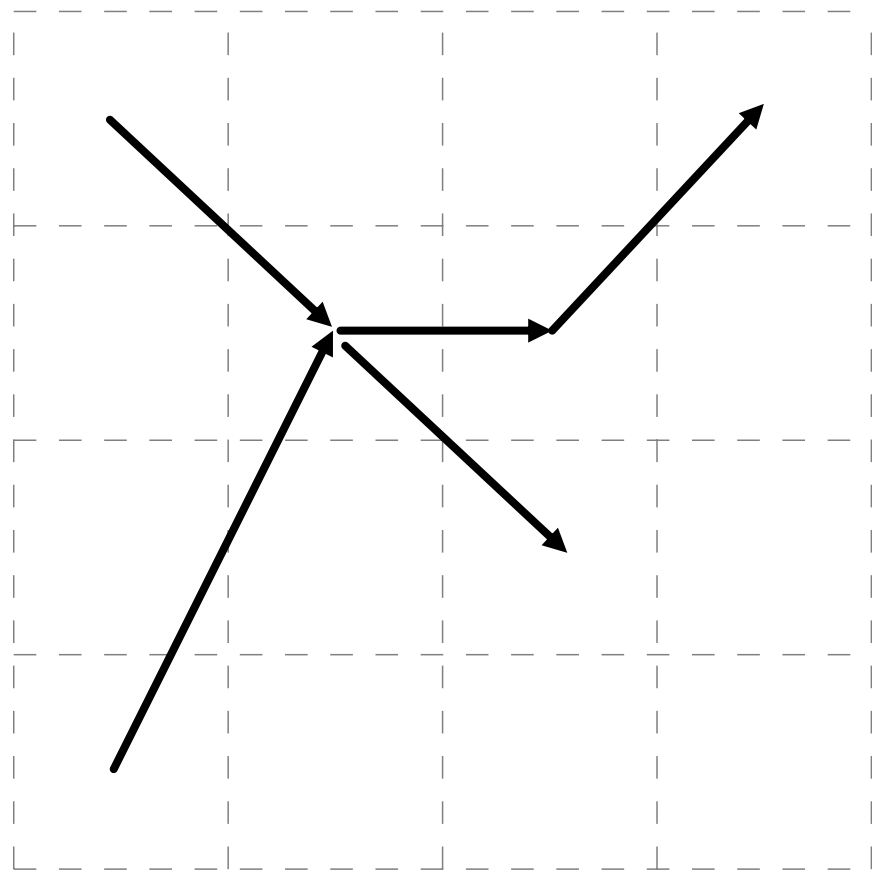


Model of Computation



Distance 6

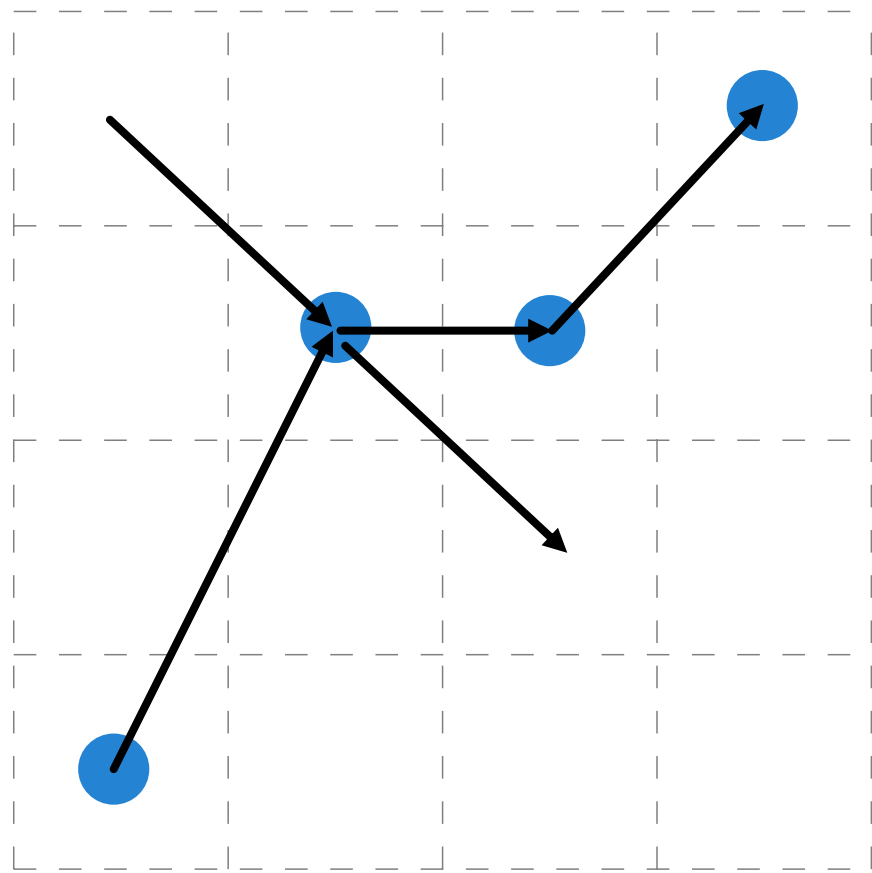
Model of Computation



Distance

Maximum 6 Total 10

Model of Computation



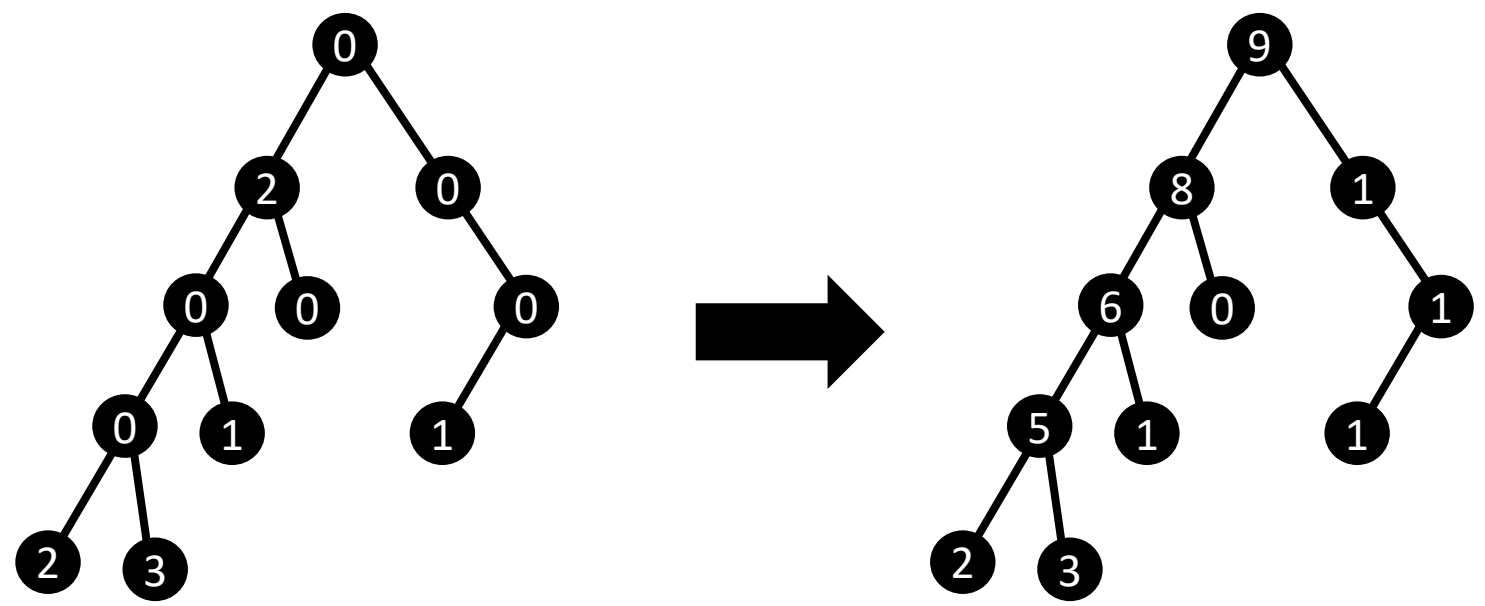
Distance

Maximum 6

Total 10

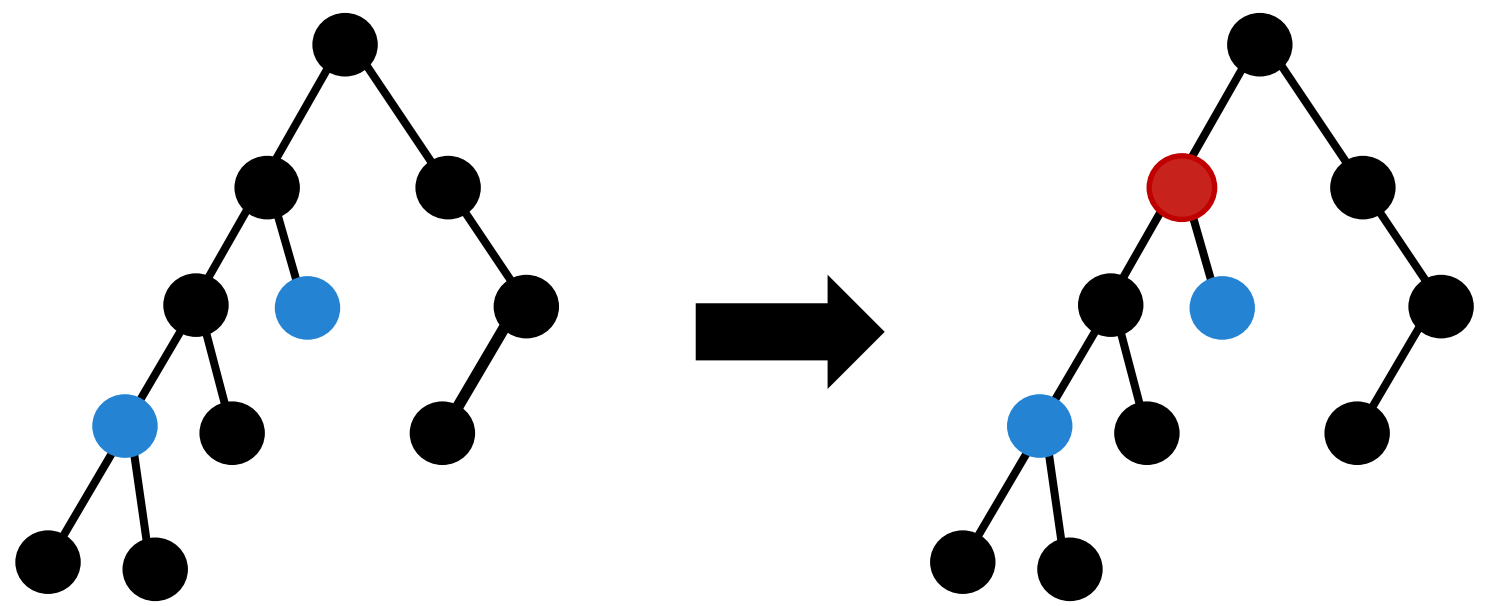
Depth 4

Problems on Trees



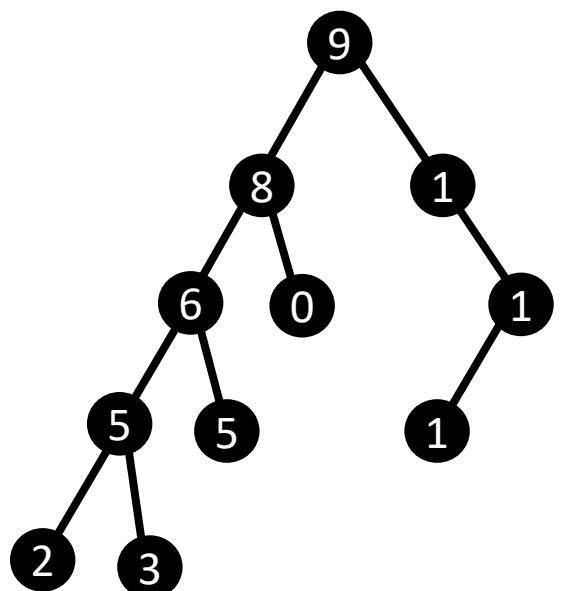
Trefix-Sum

Problems on Trees



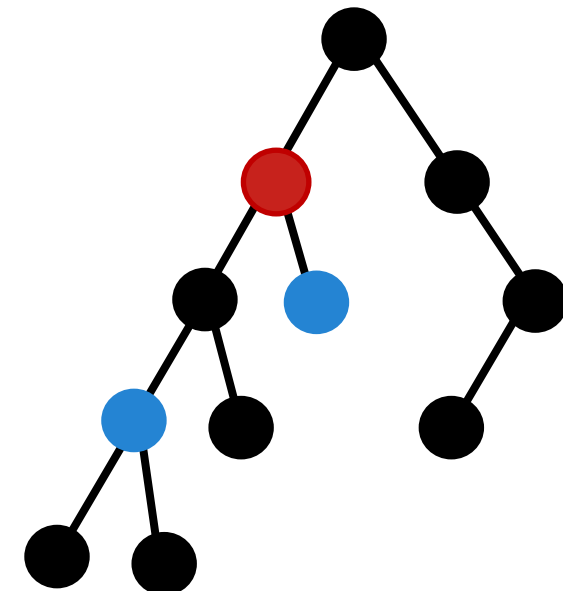
Lowest Common Ancestor (LCA)

Results



Treefix Sum
 n vertices

$O(n \log n)$ total distance
 $O(\log^2 n)$ depth



LCA
 n vertices

$O(n \log^2 n)$ total distance
 $O(\log^2 n)$ depth

Challenges

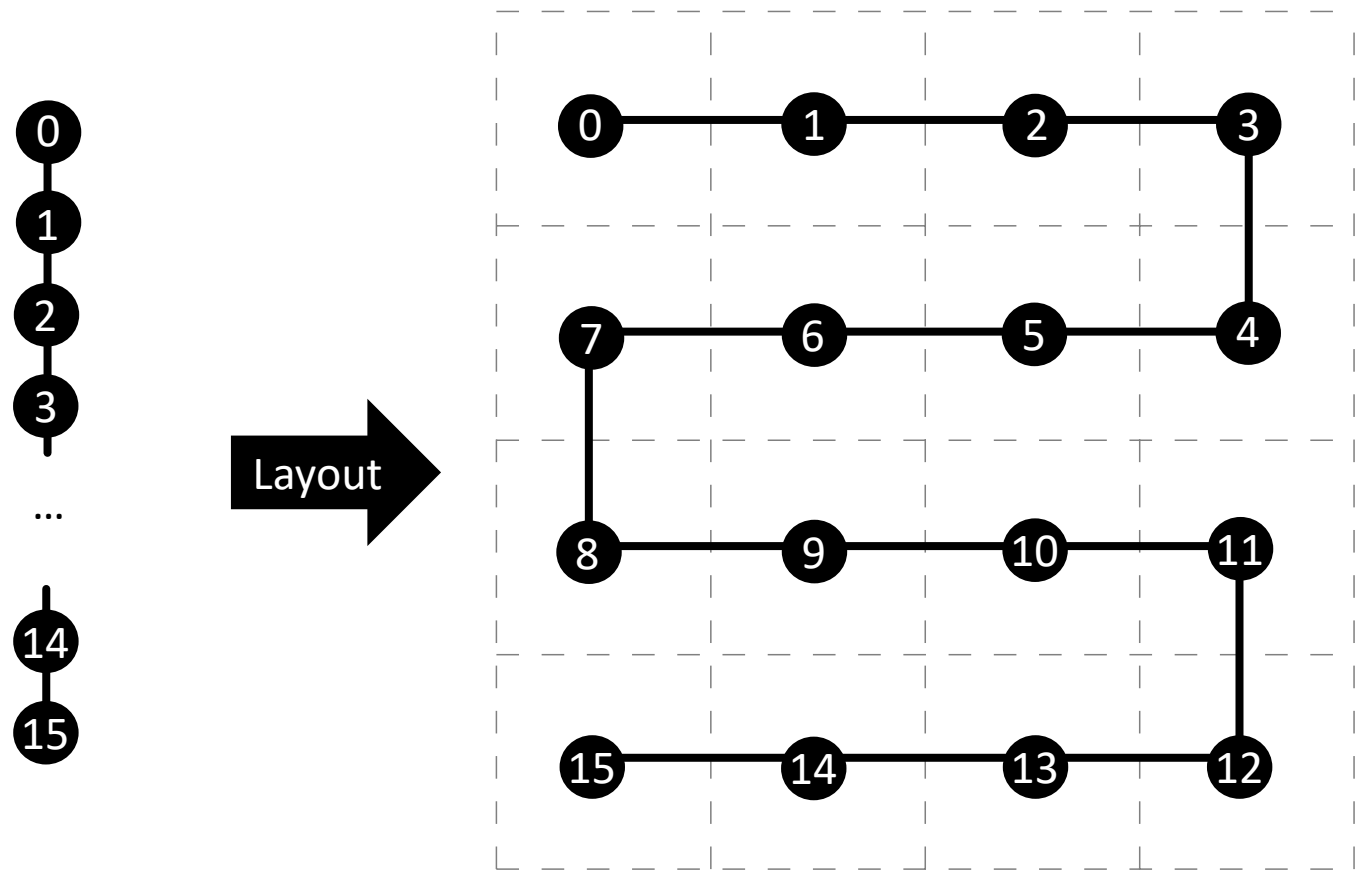
- Sparse, Irregular Structures
- Costly Random Access

PRAM Simulation:
 $\Omega(n\sqrt{n} \log n)$ total distance

Our Approach

- A. Spatial Layout
- ↓
- B. Logical Operations

Spatial Layout – Paths



Message a neighbor in the path:

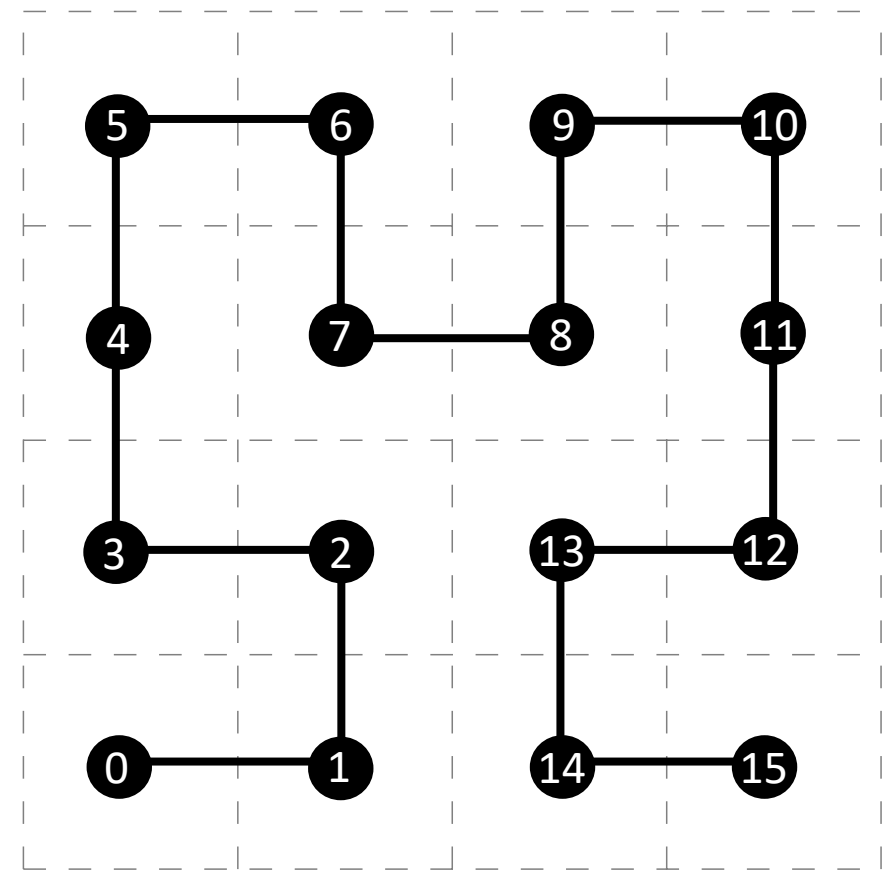
1 hop: distance 1

k hops: distance $\leq k$
average distance $\Theta(k)$

Spatial Layout – Paths



Layout →



Hilbert Curve (2D)

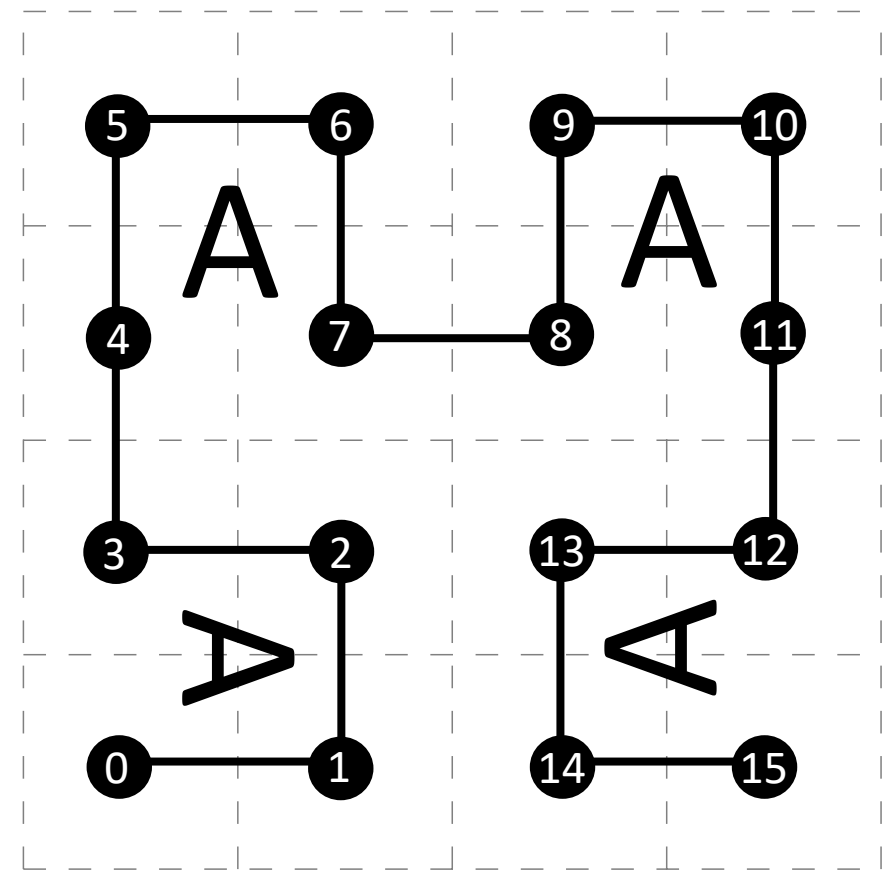
Message a neighbor in the path:

1 hop: distance 1

Spatial Layout – Paths



Layout →



Hilbert Curve (2D)

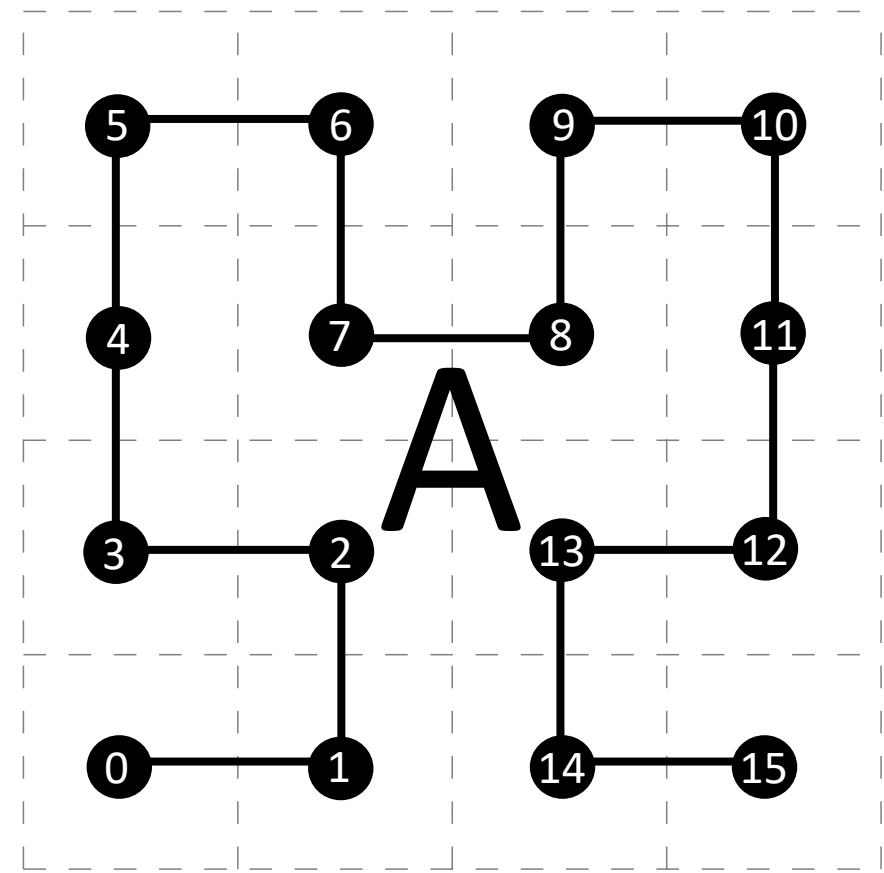
Message a neighbor in the path:

1 hop: distance 1

Spatial Layout – Paths



Layout →



Hilbert Curve (2D)

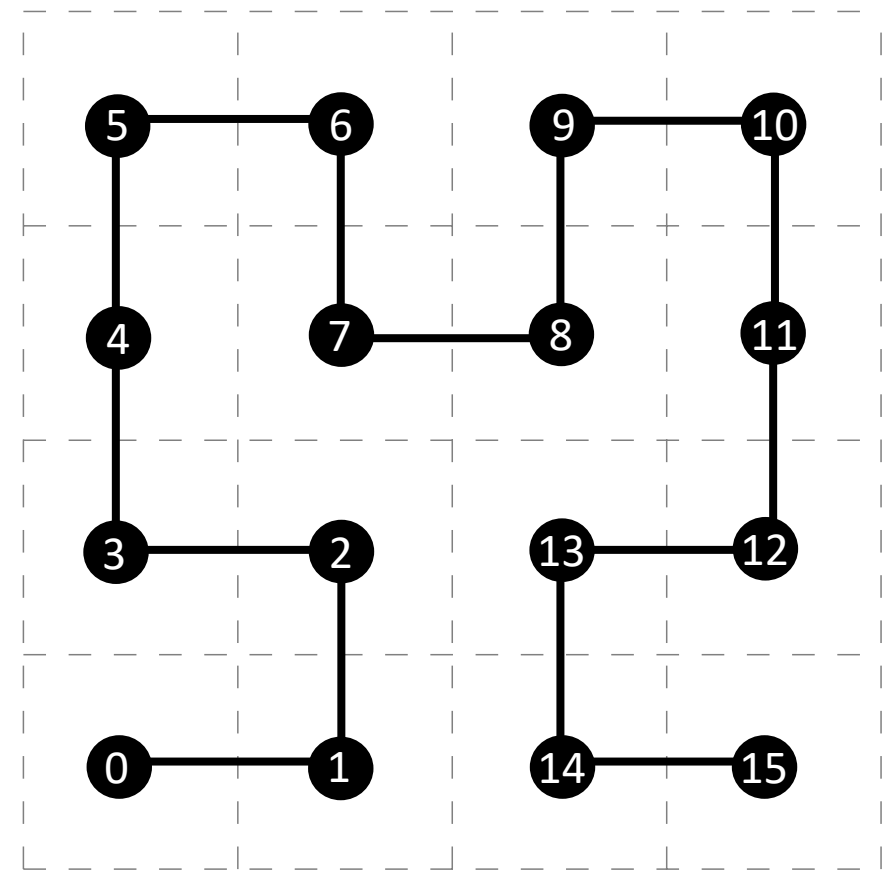
Message a neighbor in the path:

 1 hop: distance 1

Spatial Layout – Paths



Layout →



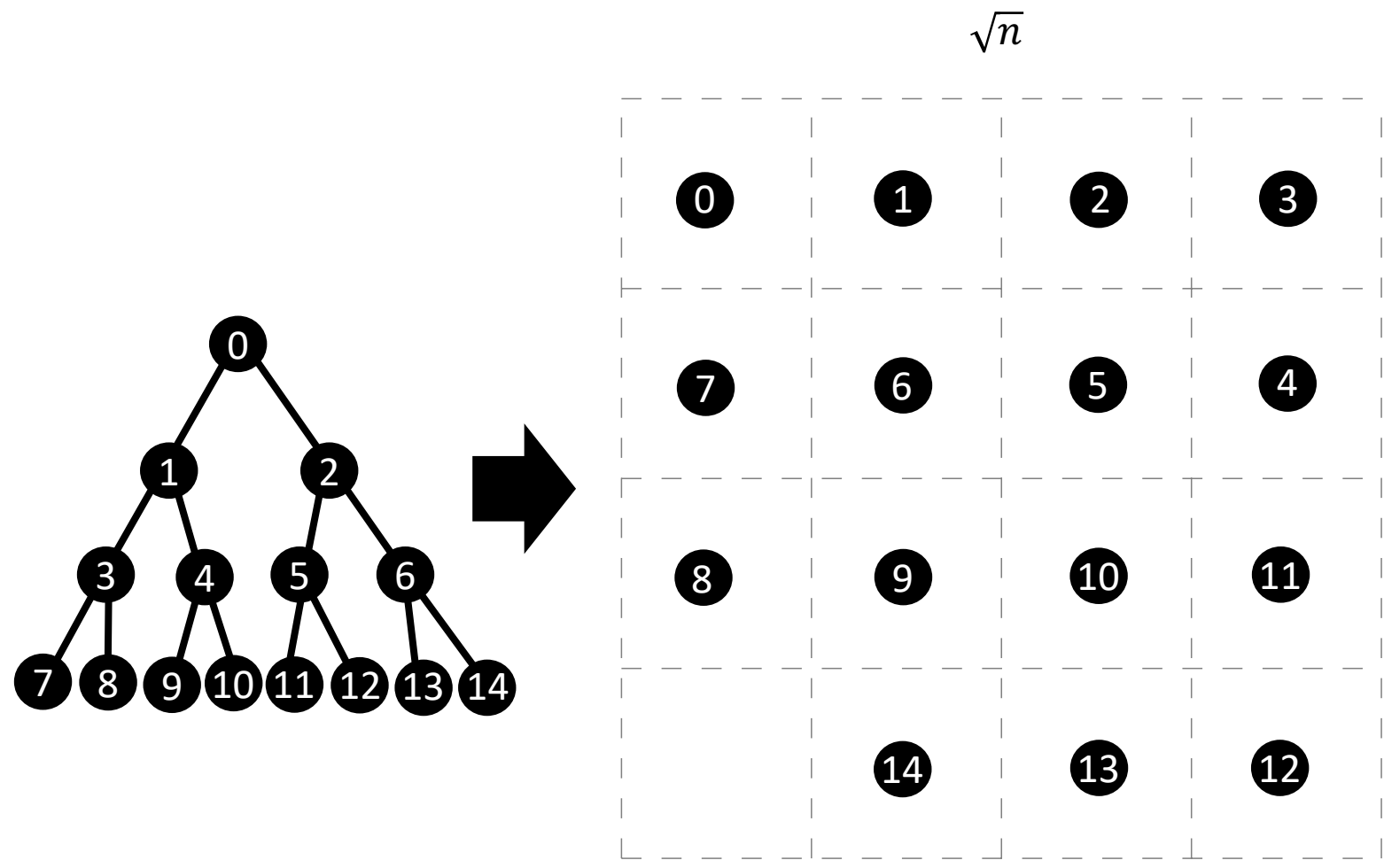
Hilbert Curve (2D)

Message a neighbor in the path:

1 hop: distance 1

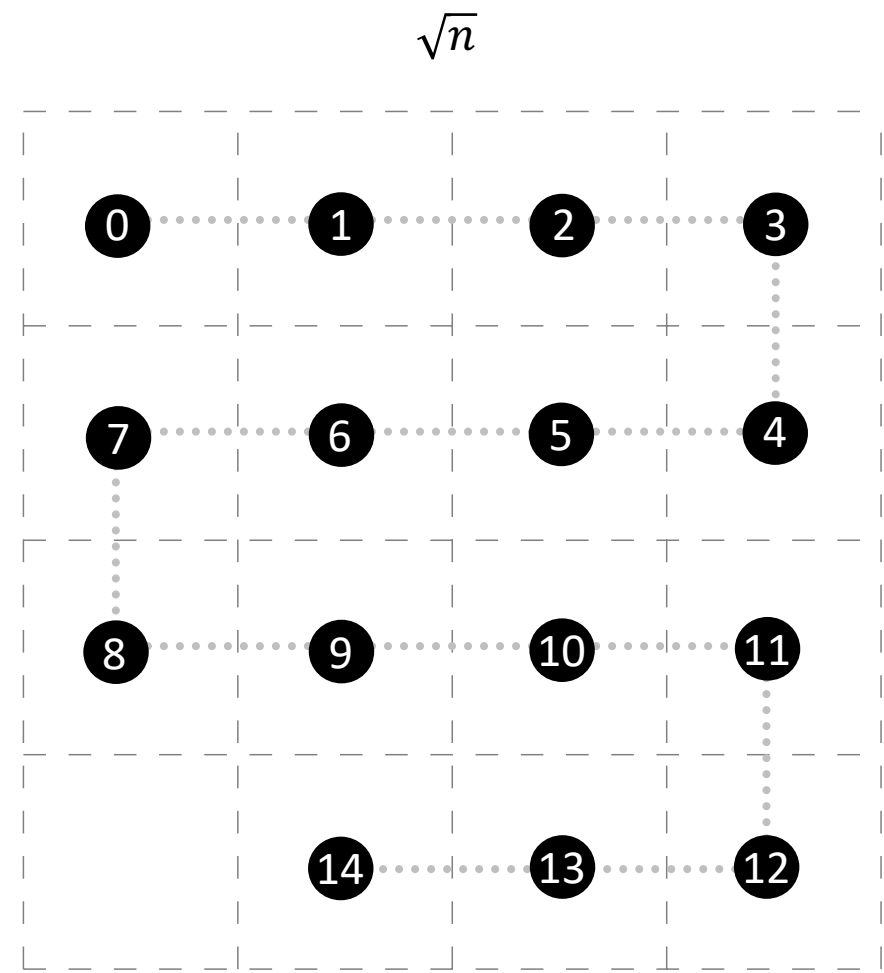
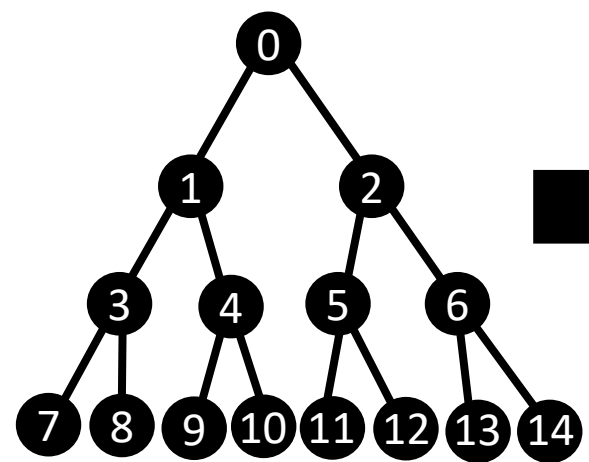
k hops: distance $\leq 3\sqrt{k} - 2$
average distance $\Theta(\sqrt{k})$

Spatial Layout – Trees



*Message a neighbor in the **tree**:*

Spatial Layout – Trees

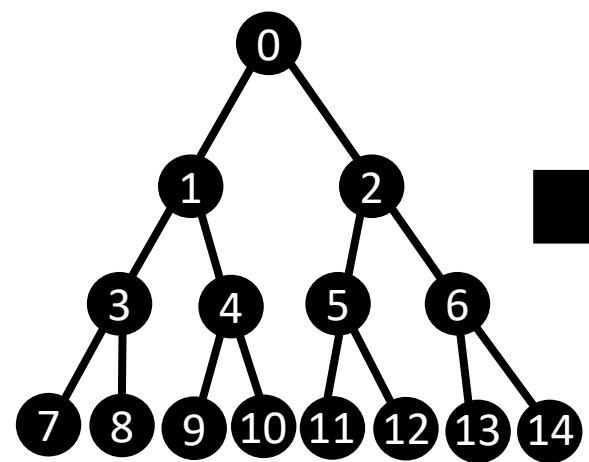


*Message a neighbor in the **tree**:*

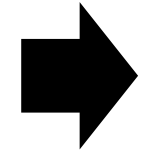
1. Linearize Tree

2. Layout According to 2D Curve

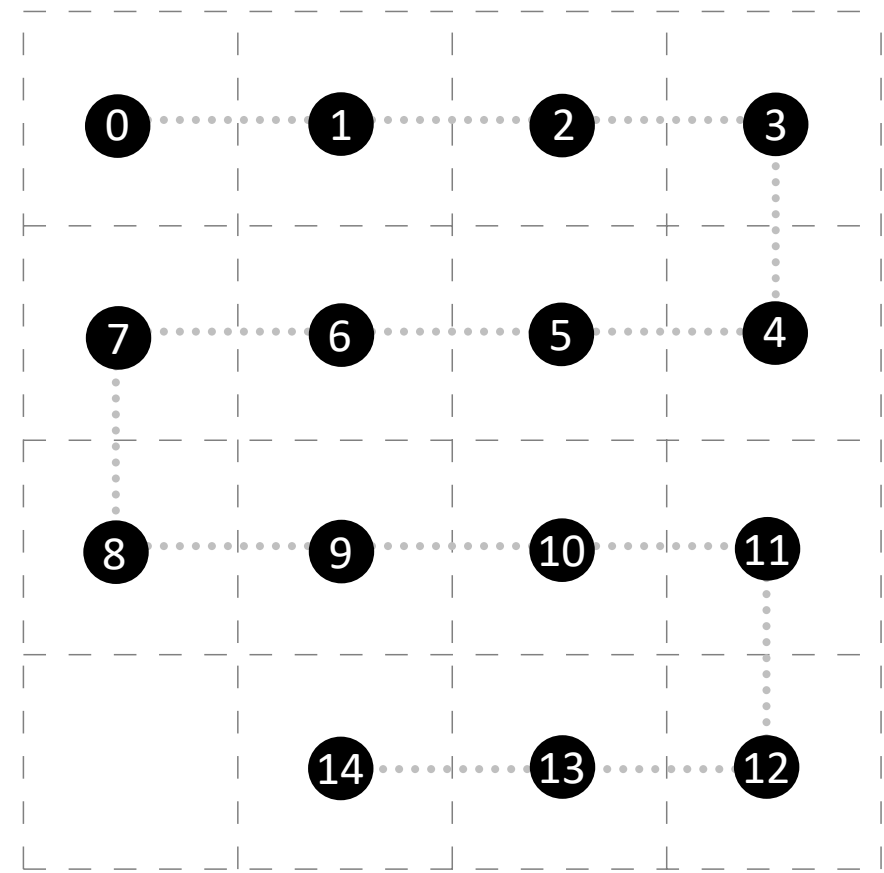
Spatial Layout – Trees



BFS



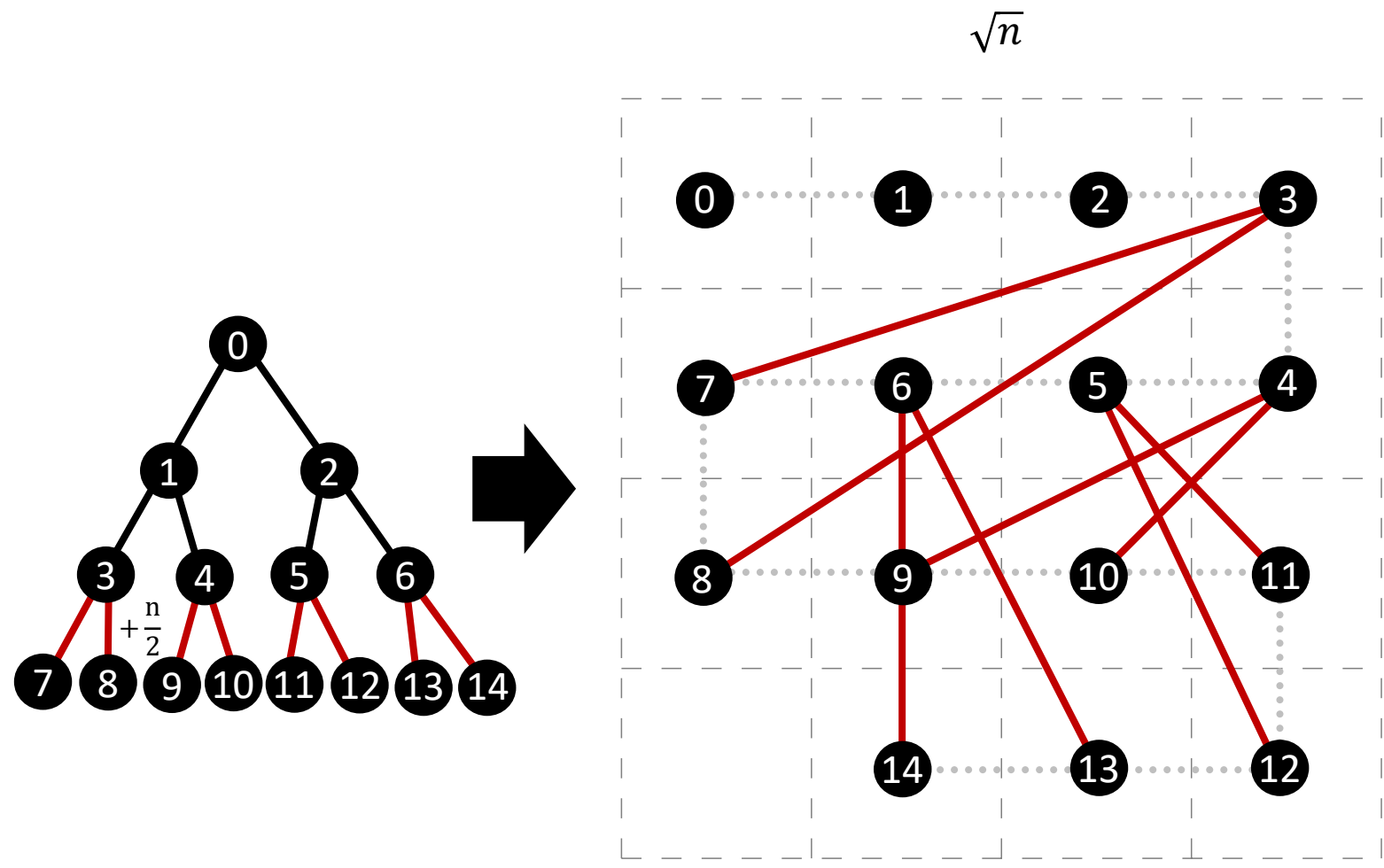
$$\sqrt{n}$$



BFS+Snake

*Message a neighbor in the **tree**:*

Spatial Layout – Trees

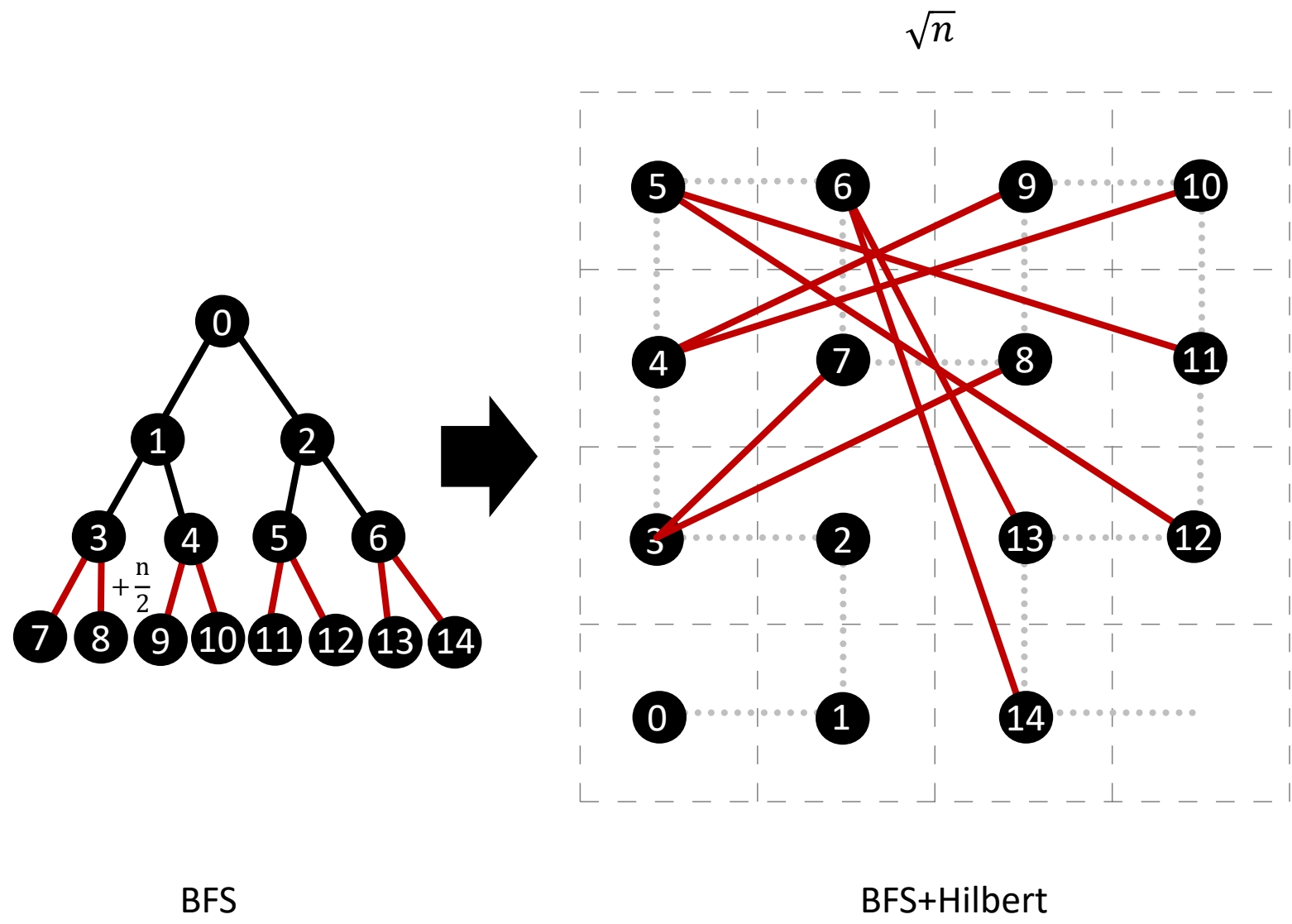


*Message a neighbor in the **tree**:*

BFS

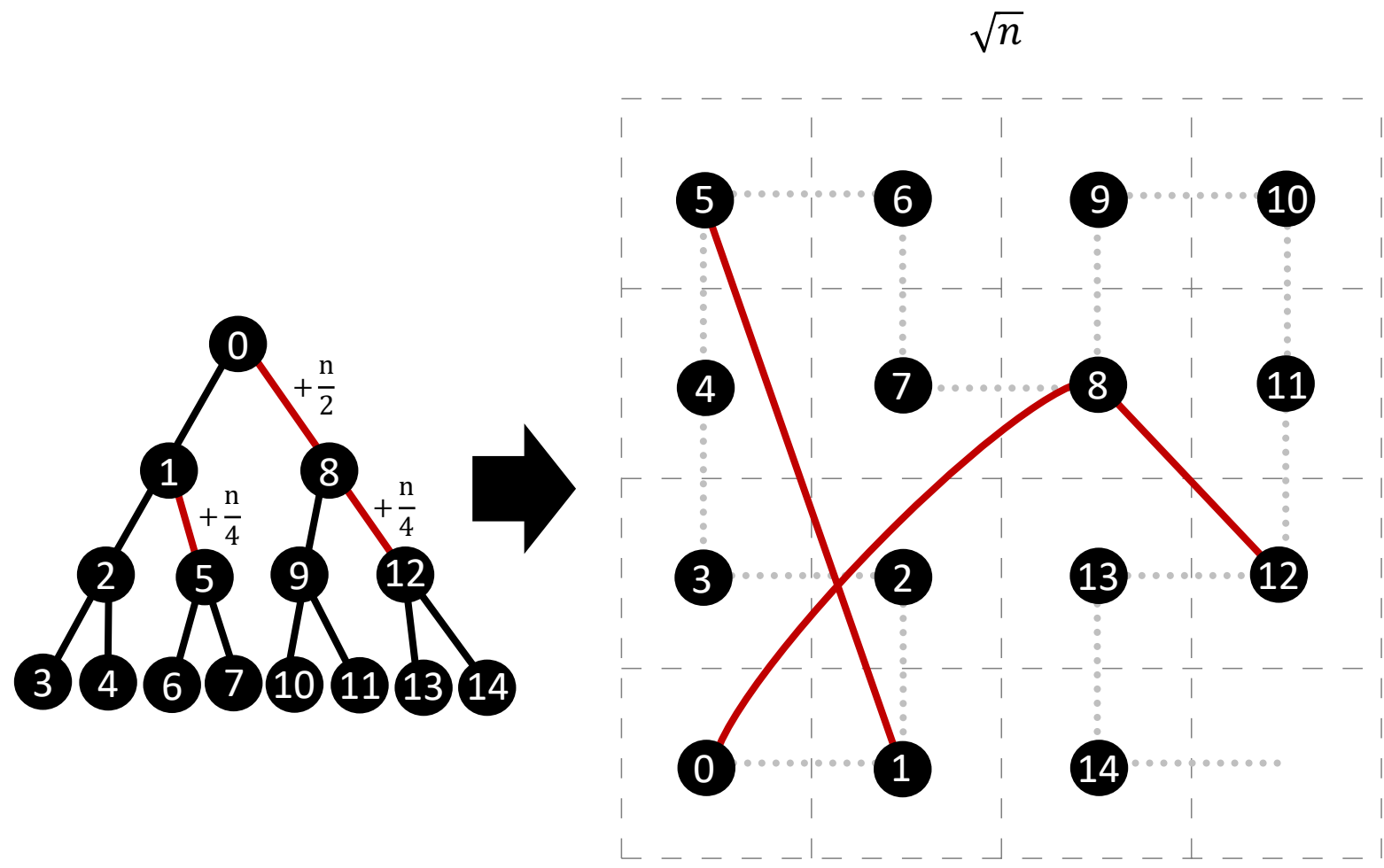
BFS+Snake

Spatial Layout – Trees



*Message a neighbor in the **tree**:*
 1 hop: average distance $\Theta(\sqrt{n})$

Spatial Layout – Trees



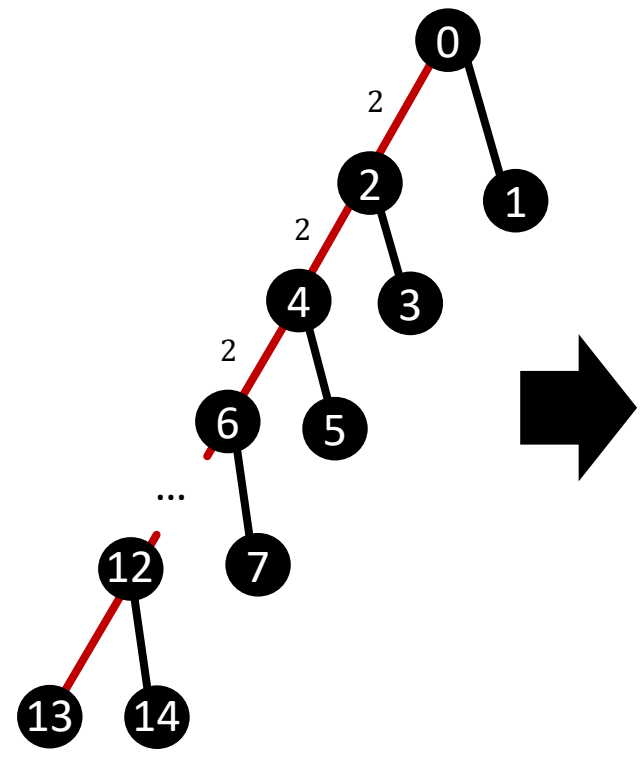
DFS

DFS+Hilbert

*Message a neighbor in the **tree**:*

1 hop: average distance $\Theta(1)$

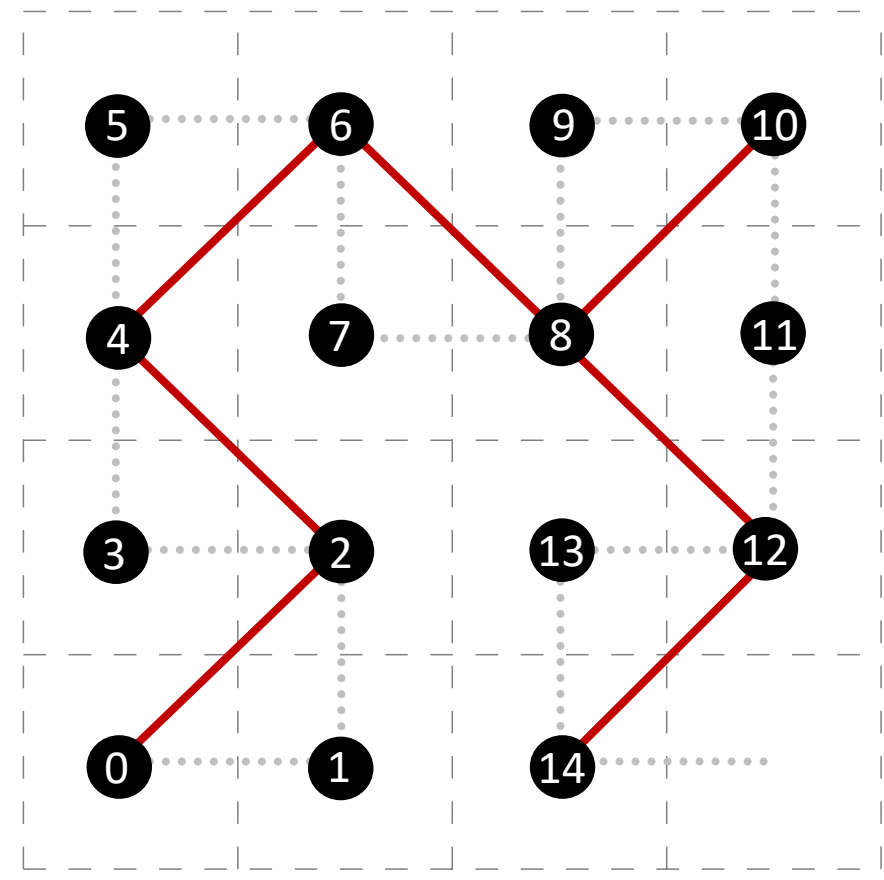
Spatial Layout – Trees



Recursively traverse children
 in order of increasing size

Light-First Order

$$\sqrt{n}$$

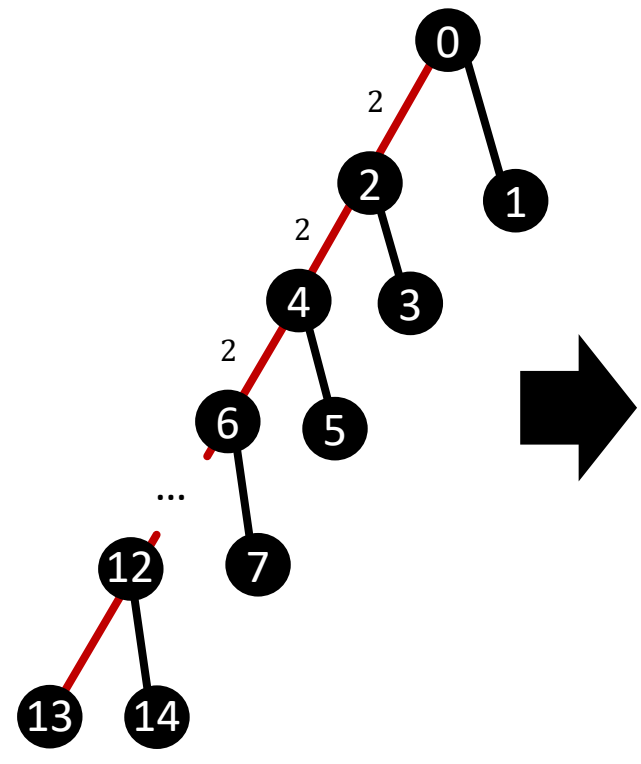


Light First+Hilbert

*Message a neighbor in the **tree**:*

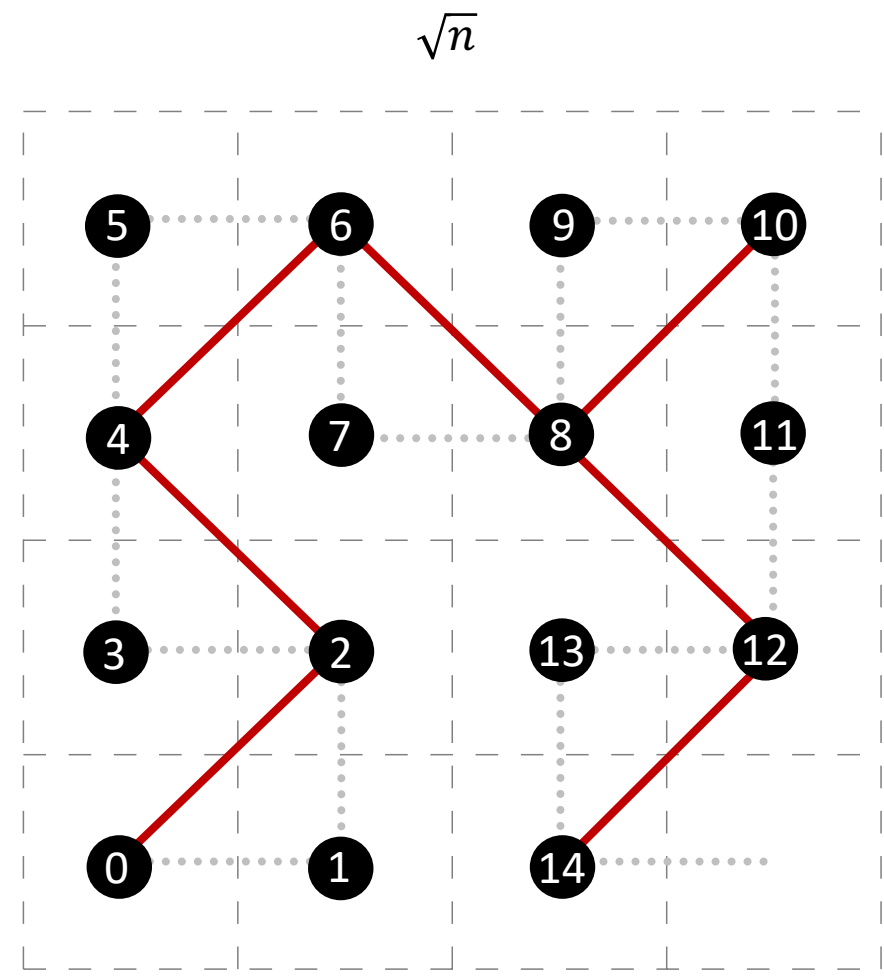
 1 hop: distance ≤ 2

Spatial Layout – Trees



Recursively traverse children in order of increasing size

Light-First Order

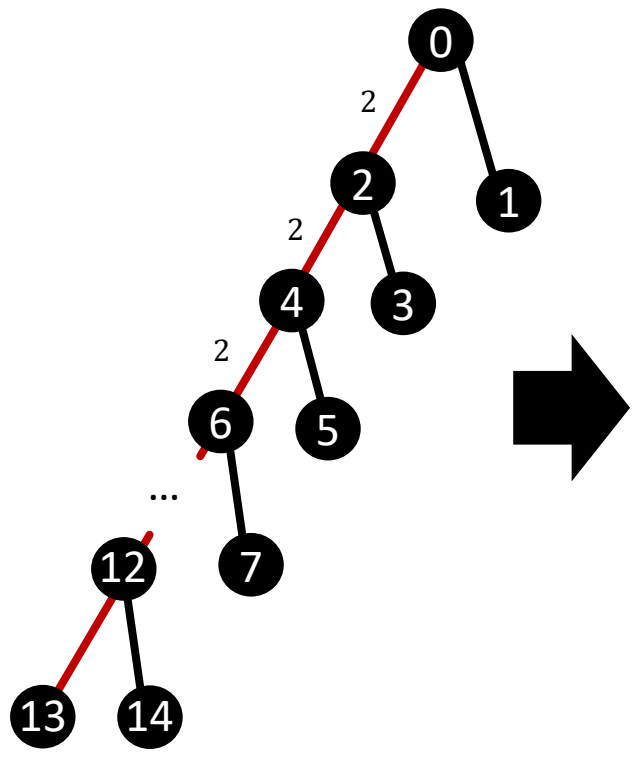


Light First+Hilbert

*Message a neighbor in the **tree**:*

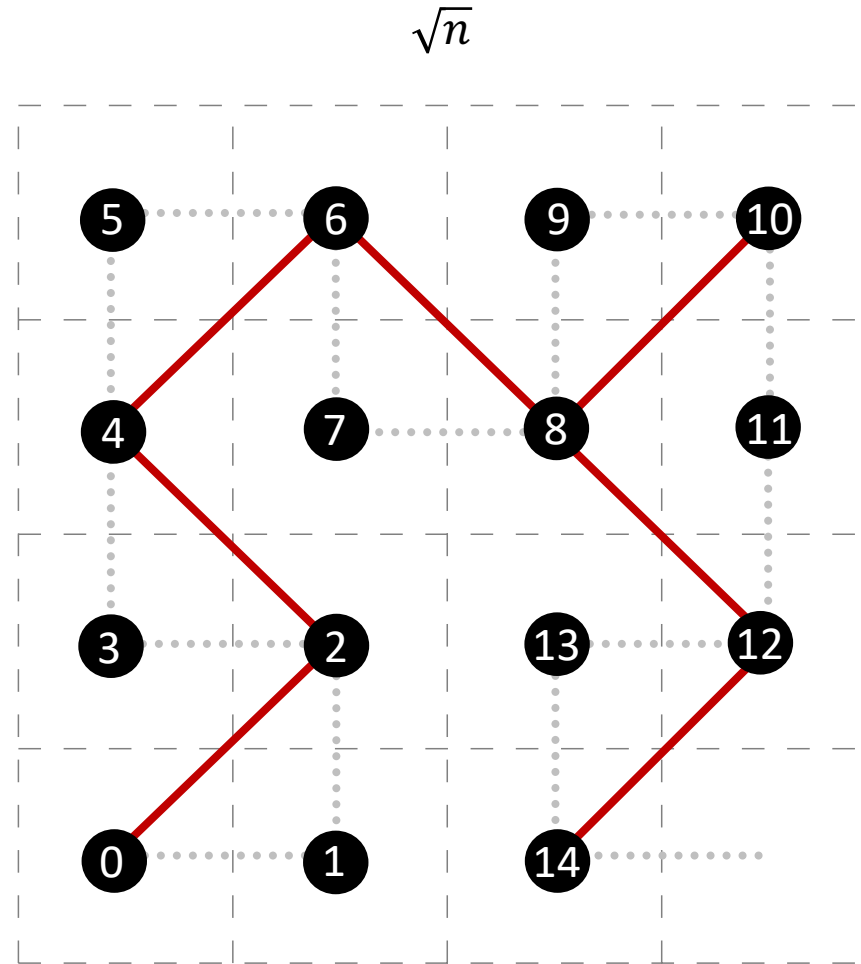
For maximum degree d ,
 Light-First + Hilbert guarantees
average distance $O(d)$

Spatial Layout – Trees



Recursively traverse children in order of increasing size

Light-First Order



Light First+Hilbert

Message a neighbor in the **tree**:

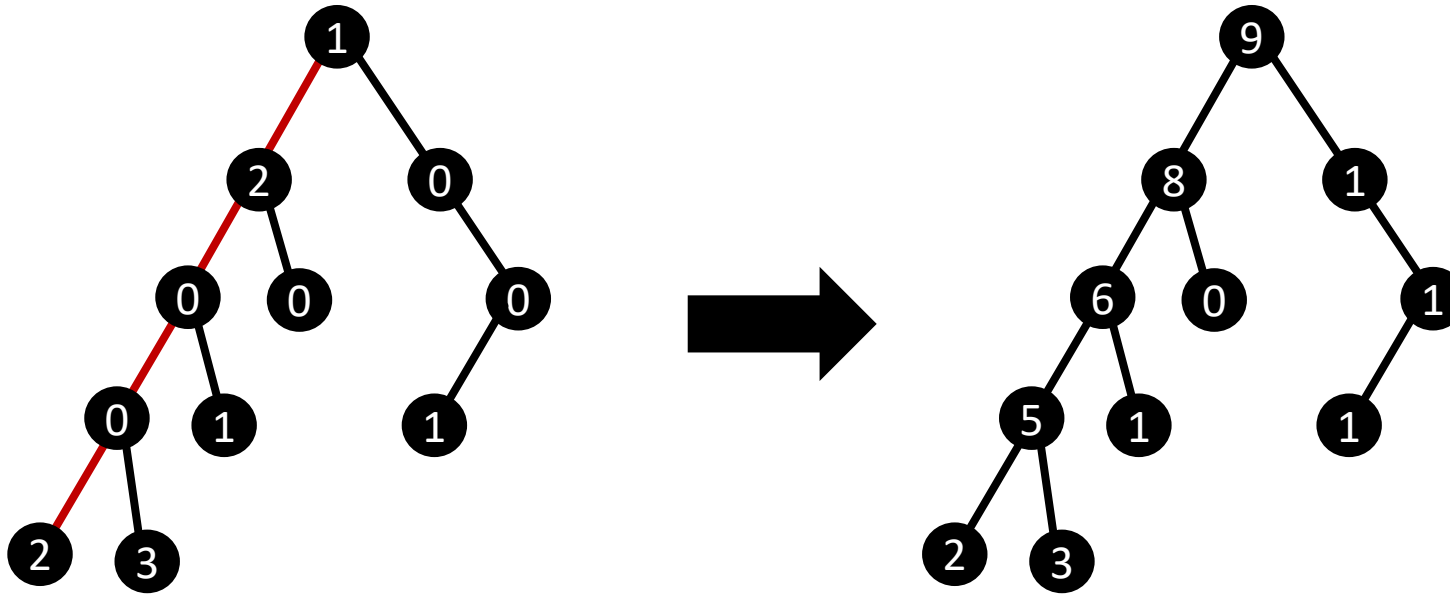
We can transform the tree into bounded degree

For maximum degree d ,
 Light-First + Hilbert guarantees
average distance $O(d)$

Optimal up to constant factors

Other space-filling curves also work

Logical Operation – Treefix Sum



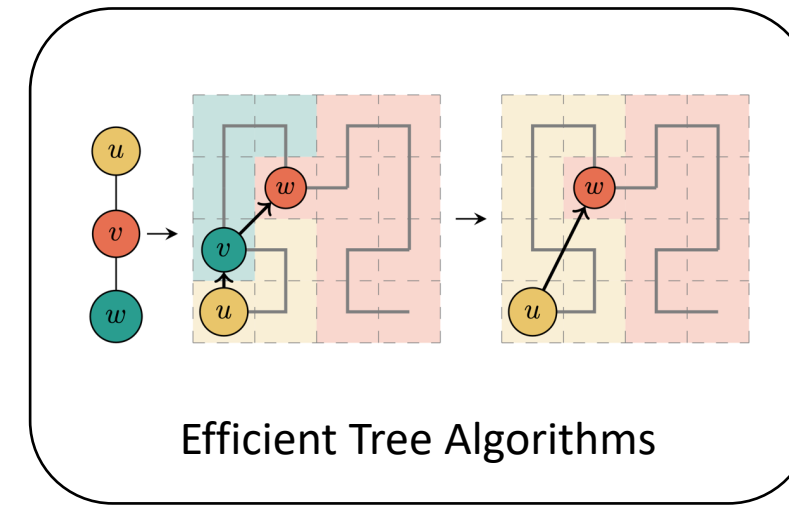
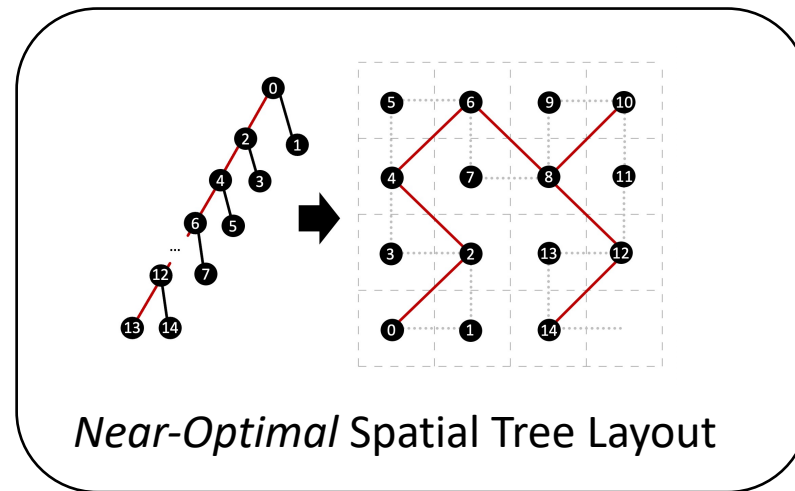
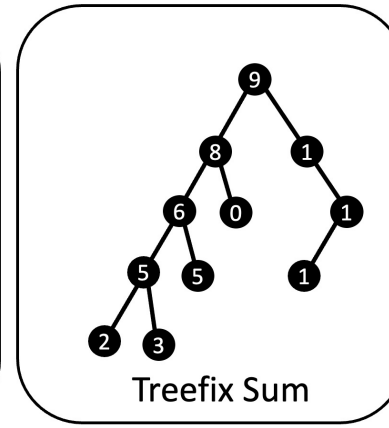
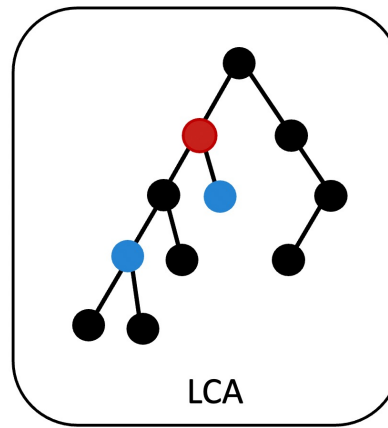
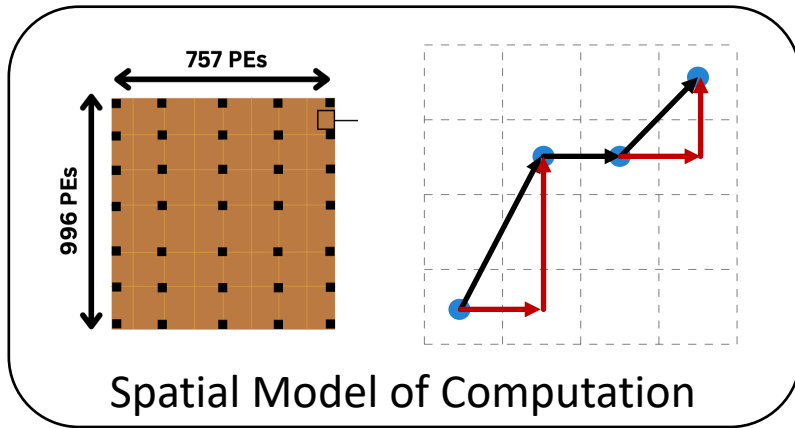
Idea:
Send partial sums to parents

Spatial layout guarantees
low average distance

Naïve Approach:
 $\Omega(n)$ depth
 n vertices

Parallel Tree Contraction:
 $O(\log^2 n)$ depth w.h.p.


Conclusions



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