Implementing Locks in C

Recitation

How to implement atomics?

Two ways:

- Use intrinsics offered by your compiler
- Use inline assembly

Today we will take a look at both options.

We use the gcc compiler as an example. Other compilers (Intel, XLC) have very similar features.

GCC Atomic Builtins

```
type __sync_fetch_and_add(*ptr, type val)
type __sync_add_and_fetch(*ptr, type val)
```

- Full memory Barrier
- Instead of "add" also available as "sub", "xor", etc.
- Returns old value (FAA) or the new one (AAF)

GCC Atomic Builtins

```
bool __sync_bool_compare_and_swap(type *ptr, type oldval, type newval)
type __sync_val_compare_and_swap(type *ptr, type oldval, type newval)
```

- Full memory Barrier
- Set *ptr to newval if *ptr is equal to oldval
- Return oldval (_val_) or truth-value of the comparison (_bool_)

GCC Atomic Builtins

```
void __sync_synchronize()
```

- Full memory barrier (mfence)

```
type __sync_lock_test_and_set(type *ptr, type
value)
```

- Full memory barrier
- Write value into *ptr, return the previous value

GCC Inline Assembly

- GCC inline Assembly uses AT&T Syntax: The first operand is the source, second operand the destination
- Intel Syntax: destination is first operand
- Register names are prefixed with %
- Constants are prefixed with $ ($0x for Hex)

Example: mov $42, %eax

GCC Inline Assembly

- The size of the operands is part of the mnemonic (optional! - assembler will guess)

```
○ movq $0x42, %reax      64 Bit
○ movl $0x42, %eax       32 Bit
○ movw $0x42, %ax        16 Bit
○ movb $0x42, %al         8 Bit
```

GCC Inline Assembly

- Memory accesses are specified as $disp(%base, %offset , $muliply) this refers to *(base + disp + offset * multiply)
- Everything except base is optional
- Common Case:
  - `movl -4(%ebp), %eax`
  - `movl (%ecx), %edx`

GCC Inline Assembly

- C compiler does not "understand" inline assembly - it simply copies it into the output stream

- We need to tell the compiler
  - Which registers we overwrite in assembly
  - Which values we access

GCC Inline Assembly

```
int a=10, b;
asm ("movl %1, %%eax;\n"
     "movl %%eax, %0;\n"
      :"=r"(b)        /* output */
      :"r"(a)         /* input */
      :"%eax"         /* clobbered register */);
```

Input/Output Values are referenced by %0, %1, etc.

Constraints tell the compiler where to put values:

"r" -> any register (a -> eax, b->ebx, etc)

"=r" -> register is used write-only

"m" -> memory location

Compare and Swap in Assembler

```
unsigned long cas(volatile unsigned long* ptr,
                  unsigned long old, unsigned long new) {
    unsigned long prev;
    asm volatile("lock; cmpxchgq %1, %2;"
                 : "=a"(prev)
                 : "r"(new), "m"(*ptr), "a"(old) : );
    return prev;
}
```

cmpxch: if (eax == dest) {ZF=1; dest=src}
else {ZF=0; eax=dest}

Fetch and Add in Assembler

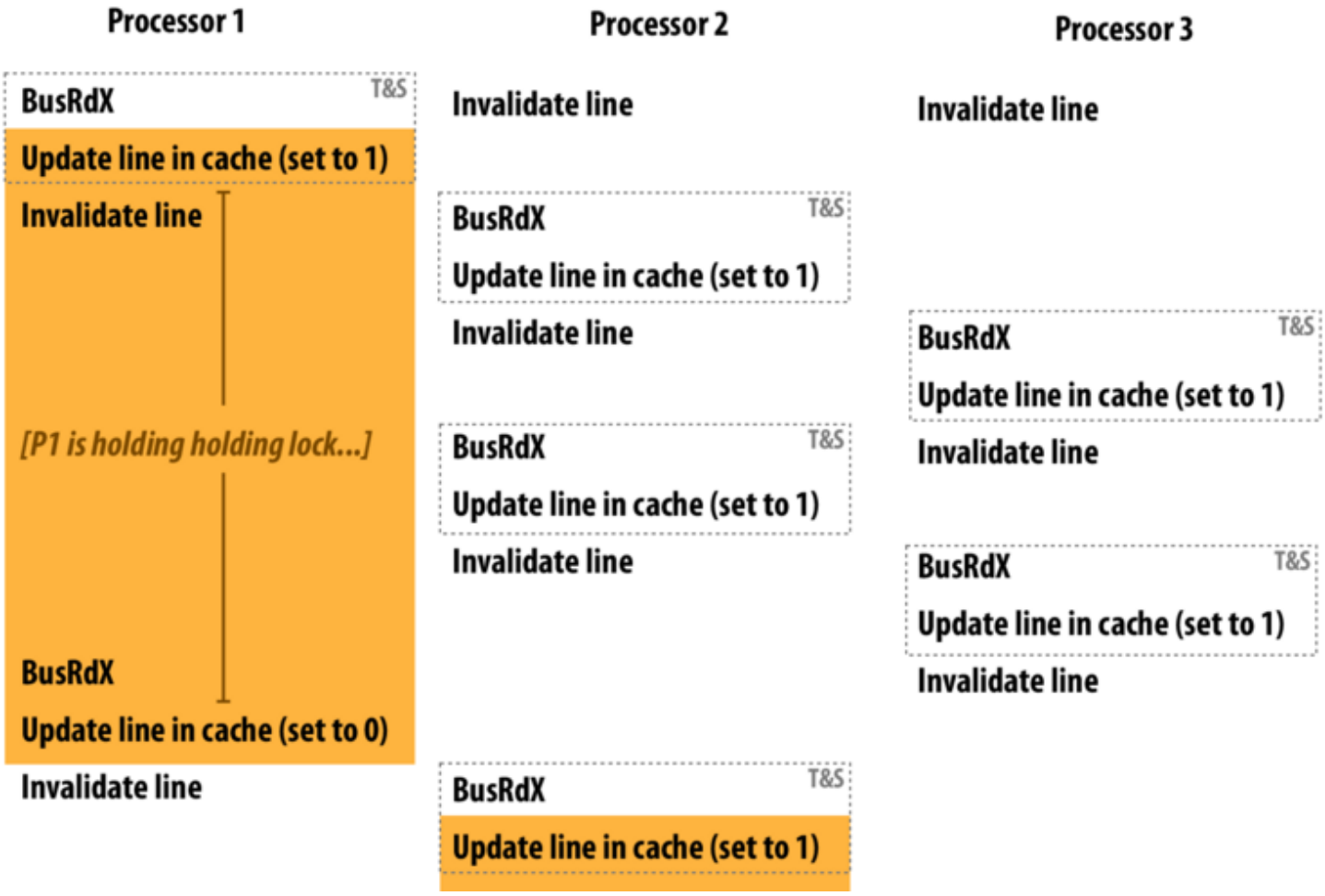
```
int fetch_and_add(int* ptr, int val){
    asm volatile(
        "lock; xaddl %%eax, %2;"
        :"=a" (val)
        : "a" (val), "m" (*ptr) : );
        return val;
 }
```

xadd: src = dest; dest += src;

Always use volatile to prevent the compiler from reordering!

Let's take a look at the Code

Test & set lock: consider coherence traffic



Test & test & set lock: coherence traffic